

CS_336
INTERACTIVE THEOREM PROVING
(Attempt 2 questions out of 3)

Question 1.

- (a) Define the notion of a *reduction system*. What does it mean for a reduction system to be *weakly normalising*? What does it mean for it to be *confluent*?

[6 marks]

- (b) Determine a reduction system which is not weakly-normalising and non-confluent, and which uses as few terms and as few reductions as possible. Explain why it is non-confluent and not weakly-normalising.

[6 marks]

- (c) Derive, using the rules for introducing terms of the simply typed λ -calculus, that

$$(\lambda f : \circ \rightarrow \circ. \lambda x : \circ. f(f\ x)) : (\circ \rightarrow \circ) \rightarrow \circ \rightarrow \circ$$

[6 marks]

- (d) Show that it is not possible to assign a simple type to $\lambda x. x\ x$, i.e. that there exist no types σ, τ such that $(\lambda x : \sigma. x\ x) : \tau$.

[7 marks]

Question 2.

- (a) Describe briefly two examples of how dependent types could be used in general programming.

[4 marks]

- (b) There are two ways of introducing the product of two sets $A :: \text{Set}$ and $B :: \text{Set}$ in Agda. Introduce both of them.

[6 marks]

- (c) Let AB be one of the products of the sets A and B , as introduced in (b). Solve the following goal in Agda:

```
f (c_a :: C → A)
  (c_b :: C → B)
  (c :: C)
  :: AB
= {!  !}
```

[6 marks]

- (d) Assuming $X, Y, Z :: \text{Set}$, introduce in Agda a function f such that

$$f :: (X \rightarrow Z) \rightarrow Z \rightarrow (X \rightarrow Y) \rightarrow (Y \rightarrow Z) \rightarrow Z$$

Hint: When we apply `intro` and then introduce one `let`-expression of type $X \rightarrow Z$, we obtain the following goals:

```
f :: (X → Z) → Z → (X → Y) → (Y → Z) → Z
= λ(x_z_z :: (X → Z) → Z) →
  λ(x_y :: X → Y) →
  λ(y_z :: Y → Z) →
  let x_z :: X → Z
    = λ(x :: X) → {!  !}
  in {!  !}
```

Solve the goals in this expression.

[9 marks]

Question 3.

- (a) Introduce in Agda the sets `Bool :: Set` of Booleans, the empty set `False :: Set`, the set `True :: Set` containing one element, and the operation `atom :: Bool → Set`, which maps a Boolean value to the formula corresponding to its truth-value.

[5 marks]

- (b) Introduce in Agda the disjoint union `X + Y :: Set` of two sets `X, Y :: Set`.

[2 marks]

In (c) - (f) assume

```
A :: Set
B :: Set
IdABool :: A → A → Bool
IdBBool :: B → B → Bool
```

The intended meaning is that `IdABool` is a Boolean valued equality on `A` and `IdBBool` is a Boolean valued equality on `B`.

- (c) Define in Agda the set theoretic equality

```
IdA :: A → A → Set
```

corresponding to `IdABool`.

[2 marks]

- (d) Introduce in Agda a set expressing “`IdA` is reflexive”.

[3 marks]

- (e) Assume that the equality `IdB` on `B` is defined as `IdA`. Introduce in Agda an equality set `IdAB` on `A + B`.

[6 marks]

- (f) Show in Agda that `IdAB` it is reflexive, provided that `IdA` and `IdB` are reflexive (where the reflexivity of `IdB` is defined as for `IdA`).

[7 marks]