

CS_376 PROGRAMMING WITH ABSTRACT DATA TYPES

(Attempt 2 questions out of 3)

Question 1.

- (a) (i) What is an abstract data type in programming? [2 marks]
- (ii) Discuss briefly the benefits of abstract data types in programming. [2 marks]
- (iii) What is the mathematical definition of an abstract data type? [2 marks]
- (iv) How are the programming and the mathematical definition of an abstract data type related? [2 marks]
- (v) What are the advantages of the mathematical definition of abstract data types over the programming definition? [2 marks]
- (b) Discuss the advantages and disadvantages of defining abstract data types in a functional and in an object-oriented programming language. What are the differences with respect to the addition of generators and observers? [6 marks]
- (c) (i) What is a *congruence* on a Σ -algebra? [2 marks]
- (ii) Let $\Sigma = (S, \Omega)$ be a signature and let $\varphi: A \rightarrow B$ be a homomorphism between Σ -algebras A and B . Let \sim be a congruence on B . For every sort $s \in S$ define a binary relation \approx_s on A_s by
- $$a \approx_s a' \quad :\iff \quad \varphi_s(a) \sim_s \varphi_s(a')$$
- Show that
- $$\approx := (\approx_s)_{s \in S}$$
- is a congruence on A . [7 marks]

Question 2.

- (a) Explain the differences between minimal, intuitionistic and classical logic. What can you say about the behavior of these logics with respect to program synthesis from proofs?

[5 marks]

- (b) (i) Derive in minimal logic the formula $((P \rightarrow Q) \rightarrow Q) \rightarrow R) \rightarrow P \rightarrow R$.

[4 marks]

- (ii) Write down the proof term corresponding to the derivation you found in (i).

[3 marks]

- (iii) Explain why the following ‘derivation’ is not correct:

$$\frac{\frac{\frac{x = x \text{ refl}}{\exists y (x = y)} \exists^+ \quad \frac{\frac{\frac{x = y}{\forall y (x = y)} \forall^+}{x = y \rightarrow \forall y (x = y)} \rightarrow^+}{\forall y (x = y \rightarrow \forall y (x = y))} \forall^+}{\forall y (x = y)} \exists^-}{\exists x \forall y (x = y)} \exists^+$$

[3 marks]

- (c) (i) State the axioms for the data type of boolean values.

[2 marks]

- (ii) Let x be a variable of sort boole. Derive from the axioms given in (i) the formula

$$\exists x ((x = T \rightarrow P) \wedge (x = F \rightarrow Q)) \rightarrow (P \vee Q)$$

in minimal logic.

[8 marks]

Question 3.

- (a) (i) Define what it means for a Σ -algebra to be *initial* in a class \mathcal{C} of Σ -algebras. [2 marks]
- (ii) What is a *loose specification*? What is a model of a loose specification? [2 marks]
- (iii) What is an *initial specification*? What is a model of an initial specification? [2 marks]
- (iv) How can a model for an initial specification be constructed? [2 marks]
- (v) The class of all models of a loose or initial specification forms an abstract data type. In which case (loose or initial) is this abstract data type monomorphic? (no proof required) [2 marks]

(b) Produce an initial specification for the algebra of *queues of natural numbers* that contains (among others) a sort `queue`, a constant `emptyqueue` for the empty queue and the following operations:

- `add`: adds a number at the end of the queue;
- `front`: returns the front element of a nonempty queue (if the queue is empty the number 0 shall be returned);
- `back`: removes the front element from a nonempty queue (if the queue is empty the empty queue shall be returned);

[7 marks]

(c) Let `nat` be a sort, `0: nat` a constant, `+, *: nat × nat → nat` operations (both used in infix notation), and `x, y, z: nat` variables. Consider the term rewriting system given by the following rules:

$$x + 0 \mapsto x$$

$$x * 0 \mapsto 0$$

$$x * (y + z) \mapsto x * y + x * z$$

Prove that this term rewriting system is terminating and confluent.

[8 marks]